## On Uplink Cloud Radio Access Networks With Interconnected Radio Units

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## Outline

- Introduction
- System Model
- Point-to-Point Compression
- Leveraging Side Information
- Joint Decompression and Decoding
- Numerical Examples
- Concluding Remarks


## Outline

## - Introduction

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## Cloud Radio Access Network (C-RAN) E]

- Base Stations (BSs) operate as radio units (RUs) [China][Simeone et al:JCN].



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## State-of-the-Art: C-RAN

- Common public radio interface (CPRI) [CPRI]
- Issued by a consortium of radio equipment manufacturers
- With the aim of standardizing the communication interface between BBU and RRHs
- Prescribes the use of
- Sampling
- Scalar quantization for the digitization of the baseband IQ samples
- 8~20 bits per I/Q sample (typically around 15 )
- Supports 3GPP GSM/EDGE, 3GPP UTRA and LTE
- Allows for star, chain, tree, ring and multi-hop fronthaul topologies
- Different bit rates up to 9.8 Gbps
- Error probability ( $10^{-12}$ ), timing accuracy ( 0.002 ppm ), delay ( $5 \mu \mathrm{~s}$ )


## State-of-the-Art: C-RAN

## RU i

Sample-wise quantization (CPRI)


## State-of-the-Art: C-RAN

## RU $i$

Sample-wise quantization (CPRI)

Time index


## RU $i$

Vector
quantization


## Source Coding Results

- Conventional source coding [EIGamal-Kim, Ch. 3]


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[ElGamal-Kim, Ch. 3] test channel: $p(\hat{Y} \mid Y)$

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$Q^{-1}()$ : Compression decoder.

An equivalent Gaussian test channel


$$
\text { with } Q \sim \mathcal{C} \mathscr{N}(0, \omega)
$$

$\left(\omega=E\left[|Q|^{2}\right]:\right.$
Quantization noise power)

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\hat{Y}=Y+Q
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\text { with } Q \sim \mathscr{C} \mathscr{N}(0, \omega)
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$\left(\omega=E\left[|Q|^{2}\right]:\right.$
Quantization noise power)

$$
\omega: I(Y ; \hat{Y}) \leq C
$$

( $I(Y ; \hat{Y})$ : Mutual information between $Y$ and $\hat{Y}$ )

## State-of-the-Art: C-RAN

- Point-to-point fronthaul compression [Hoydis et al:TSP][Zhou et al:TIT]
- Vector quantization in time domain

$$
\mathrm{CU}
$$



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## Distributed Source Coding

- Distributed source coding with side information [EIGamal-Kim, Ch. 12] test channel: $p(\hat{Y} \mid Y)$


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## Distributed Source Coding

- Distributed source coding with side information [ElGamal-Kim, Ch. 12] test channel: $p(\hat{Y} \mid Y)$ An equivalent Gaussian test channel



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- Distributed source coding with side information [ElGamal-Kim, Ch. 12] test channel: $p(\hat{Y} \mid Y)$ An equivalent Gaussian test channel

$Q():$ Compression encoder;
$\mathrm{Q}^{-1}():$ Compression decoder.

$$
\hat{Y}=Y+Q
$$

$$
\text { with } Q \sim \mathcal{C} \mathscr{N}(0, \omega)
$$

$$
\omega: I(Y ; \hat{Y} \mid Z) \leq C
$$

(Relaxed constraint than $I(Y ; \hat{Y}) \leq C)$

## State-of-the-Art: C-RAN

## - Distributed fronthaul compression

[Sanderovich et al:TIT][dCoso-Simosens][Zhou-Yu:JSAC][Park et al:SPM][Zhou et al:TIT]


## State-of-the-Art: C-RAN

## - Distributed fronthaul compression

[Sanderovich et al:TIT][dCoso-Simosens][Zhou-Yu:JSAC][Park et al:SPM][Zhou et


## State-of-the-Art: C-RAN

- Joint decompression and decoding (JDD)
[Sanderovich et al:TIT][Lim et al:TIT][Park et al:SPL][Zhou et al:TIT]
equivalent to what is latter known as: Noise Network Coding [Lim et al:TIT]


$$
\mathbf { y } _ { 1 } ^ { \mathrm { ul } } \vee \longdiv { C o m p r e s s o r }
$$



|  |
| :---: |
| $y_{2}^{\mathrm{ul}} \vee$ |
|  |
| Compressor |

## State-of-the-Art: C-RAN

- Numerical example for Wyner uplink model with $C=4$ bit/symbol



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## State-of-the-Art: Inter-RU Coop $\overline{\text { Zechnion }}$

- Inter-RU cooperation for non-cooperative cellular systems:
- Analysis for Wyner, Circular Wyner models [Simeone et al:TIT][Simeone et al:FnT]

- Other UE and/or Cell-Sites cooperation in Wyner Model [Wigger et al:TIT]


## This Work

- Inter-RU cooperation for the uplink of C-RAN:
- Analysis for circular Wyner model



## Outline

- Injtiocsuctijoss
- System Model







## System Model

- Wyner-type C-RAN uplink
- $N$ pairs of RU-UE ( $\mathscr{N}=\{1,2, \ldots, N\})$
- Fronthaul connections
- $C$ bit/symbol between RU-CU
- Uplink channel

$$
\begin{aligned}
& Y_{i}=X_{i}+\alpha X_{[i-1]}+Z_{i} \\
\text { where } & Y_{i}: \text { Rx signal RU } i, \\
& X_{i}: \text { Tx signal of UE } i, \\
& Z_{i}: \text { Noise at RU } i \text { with } Z_{i} \sim N\left(0, \sigma^{2}\right), \\
& \alpha \text { : Inter-cell channel gain with } \alpha \in[0,1] .
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<Example for $N=3$ >

## System Model

- Wyner-type C-RAN uplink
- $N$ pairs of RU-UE ( $\mathscr{N}=\{1,2, \ldots, N\})$
- Fronthaul connections
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- $B$ bit/symbol between RU-RU
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<Example for $N=3>$

## Encoding at UEs

- Encoding at UE $i$
- Message $M_{i} \in\left\{1,2, \ldots, 2^{n R_{i}}\right\}$
where $R_{i}$ is the rate of the message,
$n$ is the coding block length (assumed to be sufficiently large).


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where $R_{i}$ is the rate of the message,
$n$ is the coding block length (assumed to be sufficiently large).
- Random coding with Gaussian codebook
- Message $M_{i}$ is encoded to obtain an encoded signal

$$
X_{i} \sim N(0, P) .
$$

- Signal-to-noise ratio (SNR) of the uplink channel

$$
\mathrm{SNR}=\frac{P}{\sigma^{2}}
$$

## Oblivious Processing at RUs

- In-network processing (INP) at RU $i$

- Oblivious/Nomadic: no structure information (code-books) of UE's is available at the RUs


## Oblivious Processing at RUs

- In-network processing (INP) at RU $i$
$\hat{Y}_{B,[i-1]}=Y_{[i-1]}+Q_{B,[i-1]}$,
with quantization noise

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Q_{B,[i-1]} \sim N\left(0, \omega_{B,[i-1]}\right)
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$$
\operatorname{RU}{ }_{[i-1]}
$$

$$
\hat{Y}_{B, i}=Y_{i}+Q_{B, i}
$$

with quantization noise

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Q_{B, i} \sim N\left(0, \omega_{B, i}\right)
$$

$\mathrm{RU}{ }^{[i+1]}$




Without
side informatior

With WZ-like side information

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- In-network processing (INP) at RU $i$

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## Oblivious Processing at RUs

- In-network processing (INP) at RU $i$

$$
\hat{Y}_{C, i}=S_{i}+Q_{C, i}
$$

$$
\text { with } Q_{C, i} \sim N\left(0, \omega_{C, i}\right)
$$



- Oblivious/Nomadic: no structure information (code-books) of UE's is available at the RUs


## Decoding at CU

- Decompression and decoding at CU
- CU recovers the quantized INP output signals $\hat{Y}_{C, 1}, \hat{Y}_{C, 2}, \ldots, \hat{Y}_{C, N}$.

CU


RU 1 RU 2
$\mathrm{R} \cup N$

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- Decompression and decoding at CU
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- Then, it jointly decodes the messages $\hat{M}_{1}, \hat{M}_{2}, \ldots, \hat{M}_{N}$.

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CU


RU 1 RU 2
$\mathrm{R} \cup N$

## Decoding at CU

- Unless stated otherwise, assume that

$$
\omega_{B, i}=\omega_{B}, \omega_{C, i}=\omega_{C}, \gamma_{i}=\gamma, i \in \mathscr{N} .
$$

- Vector expression of quantized signals $\left\{\hat{Y}_{C, i}\right\}_{i \in \mathcal{N}}$
where $\mathbf{H}_{X}=\mathbf{I}+(\gamma+\alpha) \mathbf{E}_{1}+\gamma \alpha \mathbf{E}_{2}, \quad$ with $\quad \mathbf{E}_{1}=$ circulant matrix with first row $\left[\begin{array}{llll}0 & \cdots & 0 & 0\end{array}\right]$,

$$
\begin{aligned}
& \mathbf{H}_{Z}=\mathbf{I}+\gamma \mathbf{E}_{1}, \\
& \mathbf{H}_{Q}=\gamma \mathbf{E}_{1}
\end{aligned}
$$

$\mathbf{E}_{2}=$ circulant matrix with first row $\left[\begin{array}{llll}0 & \cdots & 0 & 1\end{array}\right]$.
(We have $\mathbf{E}_{1} \mathbf{E}_{1}^{T}=\mathbf{E}_{2} \mathbf{E}_{2}^{T}=\mathbf{I}, \mathbf{E}_{1} \mathbf{E}_{2}^{T}=\mathbf{E}_{1}^{T}, \mathbf{E}_{2} \mathbf{E}_{1}^{T}=\mathbf{E}_{1}$ )

## Decoding at CU

- Sum-rate $R_{\text {sum }}$ can be written as

$$
\begin{aligned}
R_{\mathrm{sum}} & =I\left(\left\{X_{i}\right\}_{i \in \mathcal{N}} ;\left\{\hat{Y}_{C, i}\right\}_{i \in \mathcal{N}}\right) \\
& =\frac{1}{2} \log _{2} \operatorname{det}\left(\mathbf{I}+P\left(\sigma^{2} \mathbf{H}_{Z} \mathbf{H}_{Z}^{T}+\omega_{B} \mathbf{H}_{Q} \mathbf{H}_{Q}^{T}+\omega_{C} \mathbf{I}\right)^{-1} \mathbf{H}_{X} \mathbf{H}_{X}^{T}\right) \\
& =\frac{1}{2} \sum_{i \in \mathcal{N}} \log _{2}\left(1+P \frac{1+(\gamma+\alpha)^{2}+\gamma^{2} \alpha^{2}+(\gamma+\alpha)(1+\gamma \alpha) \lambda_{1, i}+\gamma \alpha \lambda_{2, i}}{\sigma^{2} \gamma \lambda_{1, i}+\sigma^{2}\left(\gamma^{2}+1\right)+\omega_{B} \gamma^{2}+\omega_{C}}\right),
\end{aligned}
$$

where $\lambda_{k, l}$ : Ith largest eigenvalue of $\mathbf{E}_{k}+\mathbf{E}_{k}^{T}$ given as

$$
\lambda_{k, l}=2 \cos \left(2 k \pi \frac{l-1}{N}\right) .
$$

## Design Space

- Optimization variables
- $\omega_{B}$ : quantization noise power for RU-RU links
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- Sum-rate $R_{\text {sum }}$
- Constraints
- Capacity $B$ of RU-RU links
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- Sum-rate $R_{\text {sum }}$
- Constraints
- Capacity $B$ of RU-RU links
- Capacity $C$ of RU-CU links


Modeled differently depending on decompression strategy

## Outline

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- Siystersj Mdocjel
- Point-to-Point Compression






## Point-to-Point Compression

- In this strategy, the quantized signals $\hat{Y}_{B, i}$ and $\hat{Y}_{C, i}$ are decompressed without leveraging side information.
- Constraints on $\omega_{B}$ for RU-RU links [E|Gamal-Kim, ch. 3]

$$
I\left(Y_{i} ; \hat{Y}_{B, i}\right)=\frac{1}{2} \log _{2}\left(1+\frac{P\left(1+\alpha^{2}\right)+\sigma^{2}}{\omega_{B}}\right) \leq B .
$$

- Constraints on $\omega_{C}$ for RU-CU links [ElGamal-Kim, ch.3]

$$
I\left(S_{i} ; \hat{Y}_{C, i}\right)=\frac{1}{2} \log _{2}\left(1+\frac{\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2}}{\omega_{C}}\right) \leq C
$$

## Problem Description

- Sum-rate maximization problem (P1)

$$
\begin{array}{cl}
\underset{\omega_{B}, \omega_{C}, \gamma}{\operatorname{maximize}} & \frac{1}{2} \sum_{i \in \mathcal{N}} \log _{2}\left(1+P \frac{1+(\gamma+\alpha)^{2}+\gamma^{2} \alpha^{2}+(\gamma+\alpha)(1+\gamma \alpha) \lambda_{1, i}+\gamma \alpha \lambda_{2, i}}{\sigma^{2} \gamma \lambda_{1, i}+\sigma^{2}\left(\gamma^{2}+1\right)+\omega_{B} \gamma^{2}+\omega_{C}}\right) \\
\text { s.t. } & \frac{1}{2} \log _{2}\left(1+\frac{P\left(1+\alpha^{2}\right)+\sigma^{2}}{\omega_{B}}\right) \leq B, \\
& \frac{1}{2} \log _{2}\left(1+\frac{\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2}}{\omega_{C}}\right) \leq C .
\end{array}
$$

- Not easy to solve the problem due to the non-convexity.


## Optimization

- At optimal point, the capacity constraints should be tight.
- Without loss of optimality, we can set

$$
\begin{aligned}
& \omega_{B}=\beta_{B}\left(P\left(1+\alpha^{2}\right)+\sigma^{2}\right), \\
& \omega_{C}=\beta_{C}\left(\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2}\right),
\end{aligned}
$$

with $\beta_{B}=1 /\left(2^{2 B}-1\right)$ and $\beta_{C}=1 /\left(2^{2 C}-1\right)$.

- Therefore, the optimal value for (P1) can be found via onedimensional search over the coefficient $\gamma$.


## Outline

- Isjotocsuctijoss
- Sysitersj गuccel

- Leveraging Side Information





## Leveraging Side Information

- In this strategy, the quantized signals $\hat{Y}_{B, i}$ and $\hat{Y}_{C, i}$ are decompressed while leveraging (WZ-style) side information.
- Decompression for RU-RU links
- Uplink received signal can be leveraged as side information.
- As long as inter-cell channel gain $\alpha>0$
- Decompression for RU-CU links
- Suppose successive decompression of $\hat{Y}_{C, 1}, \hat{Y}_{C, 2}, \ldots, \hat{Y}_{C, N}$.
- At each step, previously decompressed signals can be leveraged as side information.


## Side Information for RU-RU Links

- Decompression of $\hat{Y}_{B, i}$ at RU $[i+1]$
- Leveraging side information $Y_{[i+1]}$
- Constraint on $\omega_{B}$ [ElGamal-Kim, Ch. 10]



## Side Information for RU-RU Links

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## Side Information for RU-RU Links

- Decompression of $\hat{Y}_{B, i}$ at RU $[i+1]$
- Leveraging side information $Y_{[i+1]}$
- Constraint on $\omega_{B}$ [ElGamal-Kim, Ch. 10]

$$
I\left(Y_{i} ; \hat{Y}_{B, i} \mid Y_{[i+1]}\right)=\frac{1}{2} \log _{2}\left(1+\frac{E\left[Y_{i}^{2} \mid Y_{[i+1]}\right]}{\omega_{B}}\right) \leq B, \quad \text { with } E\left[Y_{i}^{2} \mid Y_{[+1+1]}=\left(1+\alpha^{2}\right) P+\sigma^{2} .\right.
$$



## Side Information for RU-CU Links

- Decompression of $\hat{Y}_{C, 1}, \hat{Y}_{C, 2}, \ldots, \hat{Y}_{C, N}$ at CU
- Consider a successive decompression with order $\hat{Y}_{C, 1} \rightarrow \hat{Y}_{C, 2} \rightarrow \ldots \rightarrow \hat{Y}_{C, N}$



## Side Information for RU-CU Links

- Decompression of $\hat{Y}_{C, 1}, \hat{Y}_{C, 2}, \ldots, \hat{Y}_{C, N}$ at CU
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CU


## Side Information for RU-CU Links

- Decompression of $\hat{Y}_{C, 1}, \hat{Y}_{C, 2}, \ldots, \hat{Y}_{C, N}$ at CU
- Consider a successive decompression with order $\hat{Y}_{C, 1} \rightarrow \hat{Y}_{C, 2} \rightarrow \ldots \rightarrow \hat{Y}_{C, N}$
- Constraint on $\omega_{C, 1}$
- No side information when decompressing $\hat{Y}_{C, 1}$

$$
I\left(S_{1} ; \hat{Y}_{C, 1}\right)=\frac{1}{2} \log _{2}\left(1+\frac{E\left[S_{1}^{2}\right]}{\omega_{C, 1}}\right) \leq C, \quad \begin{array}{r}
\text { with } E\left[S_{1}^{2}\right] \\
=\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2} .
\end{array}
$$

## Side Information for RU-CU Links

- Decompression of $\hat{Y}_{C, 1}, \hat{Y}_{C, 2}, \ldots, \hat{Y}_{C, N}$ at CU
- Consider a successive decompression with order $\hat{Y}_{C, 1} \rightarrow \hat{Y}_{C, 2} \rightarrow \ldots \rightarrow \hat{Y}_{C, N}$
- Constraint on $\omega_{C, 1}$
- No side information when decompressing $\hat{Y}_{C, 1}$

$$
I\left(S_{1} ; \hat{Y}_{C, 1}\right)=\frac{1}{2} \log _{2}\left(1+\frac{E\left[S_{1}^{2}\right]}{\omega_{C, 1}}\right) \leq C, \quad \begin{array}{r}
\text { with } E\left[S_{1}^{2}\right] \\
=\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{b}+\left(1+\gamma^{2}\right) \sigma^{2} .
\end{array}
$$

- Constraint on $\omega_{C, i}(i>1)$
- $\hat{Y}_{C, i-1}$ is leveraged as side information when decompressing $\hat{Y}_{C, i}$.


## Problem Description

- Sum-rate maximization problem (P2)

$$
\begin{array}{cl}
\underset{\omega_{B}, \omega_{C}, \gamma}{\operatorname{maximize}} & \frac{1}{2} \sum_{i \in, \mathcal{V}} \log _{2}\left(1+P \frac{1+(\gamma+\alpha)^{2}+\gamma^{2} \alpha^{2}+(\gamma+\alpha)(1+\gamma \alpha) \lambda_{1, i}+\gamma \alpha \lambda_{2, i}}{\sigma^{2} \gamma \lambda_{1, i}+\sigma^{2}\left(\gamma^{2}+1\right)+\omega_{B} \gamma^{2}+\omega_{C}}\right) \\
\text { s.t. } & \frac{1}{2} \log _{2}\left(1+\frac{E\left[Y_{i}^{2} \mid Y_{[i+1]}\right]}{\omega_{B}}\right) \leq B, \quad i \in \mathcal{N}, \\
& I\left(S_{1} ; \hat{Y}_{C, 1}\right)=\frac{1}{2} \log _{2}\left(1+\frac{E\left[S_{1}^{2}\right]}{\omega_{C, 1}}\right) \leq C, \\
& I\left(S_{i} ; \hat{Y}_{C, i} \mid \hat{Y}_{C, i-1}\right)=\frac{1}{2} \log _{2}\left(1+\frac{E\left[S_{i}^{2} \mid \hat{Y}_{C, i-1}\right]}{\omega_{C, i}}\right) \leq C, \quad i \in \mathscr{N} \backslash\{1\} .
\end{array}
$$

- The optimization can be similarly tackled as for (P1).
$\bullet$ i.e., one-dimensional search with respect to $\gamma$.


## Optimization

- At optimal point, the capacity constraints should be tight.
- Without loss of optimality, we can set

$$
\begin{aligned}
& \omega_{B}=\beta_{B}\left(\left(1+\alpha^{2}\right) P+\sigma^{2}-\frac{\alpha^{2} P^{2}}{\left(1+\alpha^{2}\right) P+\sigma^{2}}\right), \\
& \omega_{C, 1}=\beta_{C}\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2}, \\
& \omega_{C, i}=\beta_{C}\binom{\left(\gamma_{i}^{2} \alpha^{2}+\left(\gamma_{i}+\alpha\right)^{2}+1\right) P+\gamma_{i}^{2} \omega_{B}+\left(1+\gamma_{i}^{2}\right) \sigma^{2}}{-\frac{\left[(\gamma+\alpha) P+\gamma \alpha(\gamma+\alpha) P+\gamma \sigma^{2}\right]^{2}}{\left(\gamma^{2} \alpha^{2}+(\gamma+\alpha)^{2}+1\right) P+\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2}+\omega_{C, i-1}}}, i \in \mathscr{N} \backslash\{1\},
\end{aligned}
$$

with $\beta_{B}=1 /\left(2^{2 B}-1\right)$ and $\beta_{C}=1 /\left(2^{2 C}-1\right)$.

- Therefore, the optimal value for (P2) can be found via onedimensional search over the coefficient $\gamma$.


## Outline

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- Joint Decompression and Decoding




## Joint Decompression and Decoding

- Joint decompression and decoding (JDD)
- Potentially larger rates can be achieved with JDD at CU [Sanderovich et al:TIT][Lim et al:TIT][Park et al:SPL].
- Now often seen as an instance of noisy network coding [Lim et al:TIT].
- Optimal oblivious processing [Aguerri et al:arXiv]


## CU



## Joint Decompression and Decoding

- Joint decompression and decoding (JDD)
- Achievable sum-rate under JDD for given $\omega_{B}, \omega_{C}, \gamma$ [Sanderovich et al:TIT] [Lim et al:TIT]

$$
\begin{aligned}
R_{\text {sum }} & =\min _{\mathcal{S} \subseteq \mathcal{N}}\left\{|\mathcal{S}| C-\sum_{i \in \mathcal{S}} I\left(S_{i} ; \hat{Y}_{C, i} \mid \mathbf{X}\right)+I\left(\mathbf{X} ;\left\{\hat{Y}_{C, i}\right\}_{i \in \mathcal{N} \backslash S}\right)\right\} \\
& =\min _{\mathcal{S} \subseteq \mathcal{N}}\left\{|\mathcal{S}|\left(C-g_{C}\left(\omega_{B}, \omega_{C}, \gamma\right)\right)+f_{C, \mathcal{S}}\left(\omega_{B}, \omega_{C}, \gamma\right)\right\}
\end{aligned}
$$

where $g_{C}\left(\omega_{B}, \omega_{C}, \gamma\right)=\frac{1}{2} \log _{2}\left(1+\frac{\gamma^{2} \omega_{B}+\left(1+\gamma^{2}\right) \sigma^{2}}{\omega_{C}}\right)$,

$$
\begin{aligned}
& f_{C, s}\left(\omega_{B}, \omega_{c}, \gamma\right)=\frac{1}{2} \log _{2} \operatorname{det}\left(\mathbf{I}+P\left(\sigma^{2} \mathbf{H}_{Z, s} \mathbf{H}_{Z, s}^{T}+\omega_{B} \mathbf{H}_{Q, s} \mathbf{H}_{Q, s}^{T}+\omega_{C} \mathbf{I}\right)^{-1} \mathbf{H}_{X, S} \mathbf{H}_{X, S}^{T}\right), \\
& \mathbf{H}_{X, s}, \mathbf{H}_{Z, s}, \mathbf{H}_{Q, s}: \text { Submatrices of } \mathbf{H}_{X}, \mathbf{H}_{Z}, \mathbf{H}_{Q} \text { with rows in } \mathcal{S} \text { removed. }
\end{aligned}
$$

## Problem Description

- Sum-rate maximization problem (P3)

$$
\begin{array}{cl}
\underset{\omega_{B}, \omega_{C}, \gamma, R_{\mathrm{sum}}}{\operatorname{s.t.}} & R_{\text {sum }} \\
& R_{\text {sum }} \leq|\mathcal{S}|\left(C-\tilde{g}_{C}\left(\omega_{B}, \omega_{C}, \gamma\right)\right)+f_{C, \mathcal{S}}\left(\omega_{B}, \omega_{C}, \gamma\right), \quad \mathcal{S} \subseteq \mathscr{N}, \\
& \frac{1}{2} \log _{2}\left(1+\frac{E\left[Y_{i}^{2} \mid Y_{[i+1]}\right]}{\omega_{B}}\right) \leq B, \quad i \in \mathscr{N},
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& \frac{1}{2} \log _{2}\left(1+\frac{E\left[Y_{i}^{2} \mid Y_{[i+1]}\right]}{\omega_{B}}\right) \leq B, \quad i \in \mathscr{N},
\end{array}
$$

- We propose to perform one-dimensional search w.r.t. $\gamma$.
- For given $\gamma$, optimizing $\omega_{B}$ and $\omega_{C}$ is a difference-of-convex (DC) problem.
- Thus, suboptimal solution of $\omega_{B}$ and $\omega_{C}$ for given $\gamma$ can be found via concave convex procedure (CCCP) approach.


## Outline

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- Numerical Examples
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## Cut-Set Upper Bound

- For reference, we consider the Cut-Set upper bound on $R_{\text {sum }}$ as

$$
R_{\text {sum }} \leq \min \left\{N C, R_{\text {full }}\right\},
$$

where $R_{\text {full }}$ is the sum-rate achievable when full cooperation among RUs is possible, i.e.,

$$
\begin{aligned}
R_{\text {full }} & =I\left(\left\{X_{i}\right\}_{i \in \mathcal{N}} ;\left\{Y_{i}\right\}_{i \in \mathscr{N}}\right) \\
= & \frac{1}{2} \log _{2} \operatorname{det}\left(\mathbf{I}+P\left(\sigma^{2} \mathbf{H}_{Z} \mathbf{H}_{Z}^{T}\right)^{-1} \mathbf{H}_{X} \mathbf{H}_{X}^{T}\right) .
\end{aligned}
$$

## Oblvious Upper Bound

- We also consider an oblivious upper bound.
- Sum-rate that can be achieved when the RUs are co-located and send jointly quantized signals of $\left\{Y_{i}\right\}_{i \in \mathcal{N}}$ to the CU.



Enabling full RU cooperation


Achievable rate was analyzed
in [dCoso-Simoens, Thm. 1].

## Numerical Example

- Per-UE rate versus RU-RU capacity $B$
- $N=3, \mathrm{SNR}=20 \mathrm{~dB}, \alpha=0.7$

- With INP, the performance approaches upper bound as $B$ increases.


## Numerical Example

## - Per-UE rate versus RU-RU capacity $B$

- $N=3, \mathrm{SNR}=20 \mathrm{~dB}, \alpha=0.7$

- Wititn INP, trine performance approaches upper bourid as B increases.
- Leveraging SI for RU-RU link provides a slight sum-rate gain.


## Numerical Example

## - Per-UE rate versus RU-RU capacity $B$

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## - Per-UE rate versus RU-RU capacity $B$

- $N=3, \mathrm{SNR}=20 \mathrm{~dB}, \alpha=0.7$

- Witín INP, trne performance approaches upper 'Dourid as B ifucreases.
- Leveragirig SI for RIU-FiU links provides a slightit surn-rate gain.
- Leveragirng Sl for RUU-CU link leads to a significant surn-rate gain especially for srrall B.
- JDD further improves the sum-rate performance.
(This is the optimal oblivious processing [Aguerri et al:arXiv].)


## Numerical Example

- Per-UE rate versus SNR $P / \sigma^{2}$
- $N=3, \alpha=0.7, C=B \in\{1,2\}$

- In low-to-intermediate SNR regime, the gap to cutset upper bound is still large.


## Numerical Example

- Per-UE rate versus SNR $P / \sigma^{2}$
- $N=3, \alpha=0.7, C=B \in\{1,2\}$


In low-to-intermediate SNR regime, the gap to cutset upper bound is still large.

- This calls for the development of
- Improved scheme based on
- Non-oblivious RU processing
- Improved upper bound
- Extending the idea as [Wu et al:arXiv]


## Outline

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- Concluding Remarks


## Concluding Remarks

- We have studied the role of inter-RU links for improving the sum-rate of C-RAN uplink.
- Under the assumptions of
- Oblivious processing at RUs
- Wyner-type Gaussian channel


## Concluding Remarks

- We have studied the role of inter-RU links for improving the sum-rate of C-RAN uplink.
- Under the assumptions of
- Oblivious processing at RUs
- Wyner-type Gaussian channel
- Future work
- Possible optimality of non-oblivious processing also for interconnected radio units, ala:
[Aguerri, Zaidi, Caire and Shamai arXiv:1701.07237, Jan. 2017]
- Non-oblivious processing at RUs
- Compute-and-Forward based techniques [Aguerri-Zaidi][Hong-Caire]
- Edge processing
- Improved outer bounds over the cut-set bound, extending ideas as: [Wu et al:arXiv][Bidokhti et al, ISIT2017]


## Concluding Remarks

- Future work (ctd')
- C-RAN uplink set-ups with fading channels
- Downlink of C-RAN (Oblivious and Non-oblivious schemes) [Wang et al IT, Aug 2018]
- Possibly with edge processing or edge caching


## Thank you!

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## Abstract

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"On Uplink Cloud Radio Access Networks With Interconnected Radio Units,"
We address the achievable sum-rate for the cloud radio access network (C-RAN) uplink operating in a linear Wyner-type topology. In the system, a set of radio units (RUs) is connected to a control unit (CU) by means of digital finite-capacity fronthaul links, and the messages sent by the users equipment (UEs) served by the RUs are jointly decoded at the CU based on the compressed baseband signals received on the fronthaul links. The potential advantages of utilizing the inter-RU links to improve the sum-rate performance is examined. In the considered strategy, each RU performs in-network processing of the uplink received signal and of the compressed baseband signal received from the adjacent RU, with the CU performing channel decoding incorporating the in-network processing output signals. A closed-form expression of the achievable sum-rate is derived assuming point-to-point compression, and analytic expressions for other advanced compression options, leveraging side information are also provided. Insights into the advantages of inter-RU communications follow some numerical examples highlighting the performance gap to the associated sumrate upper bounds.

Joint work with Seok-Hwan Park (Chonbuk National University, Korea) and Osvaldo Simeone (King's College London).
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